

# Rectangularity theorem for conservative algebras

Libor Barto

e-mail: libor.barto@gmail.com

Charles University in Prague, Czech Republic

The Constraint Satisfaction Problem over  $\mathbb{A} = (A, \Gamma)$ , where  $A$  is a finite set and  $\Gamma$  is a set of relations on  $A$  (sometimes assumed to be finite), is a decision problem asking whether an input positive primitive (pp-) sentence over  $\mathbb{A}$  is true. A pp-sentence is a sentence formed from relations in  $\Gamma$  using conjunction and existential quantification, i.e. a sentence of the form

$$(\exists x_1 \in A) (\exists x_2 \in A) \dots R_1(\text{variables}) \wedge R_2(\text{vars}) \wedge \dots \wedge R_k(\text{vars}),$$

where  $R_i \in \Gamma$ . This decision problem is denoted by  $\text{CSP}(\mathbb{A})$ . Many well known problems, such as satisfiability problems for boolean formulas (3-SAT, HORN-SAT, ...), colorings of graphs, solving systems of equations over finite domains, etc. are examples of CSPs over suitably chosen  $\mathbb{A}$ .

The central theoretical problem in this area is the Dichotomy Conjecture of Feder and Vardi.

**Conjecture 1** (CSP Dichotomy Conjecture). *For every  $\mathbb{A}$ , the problem  $\text{CSP}(\mathbb{A})$  is either poly-time, or NP-complete.*

The algebraic approach to the Dichotomy Conjecture is based on the fact that the complexity of  $\text{CSP}(\mathbb{A})$  is determined by the clone  $\mathbf{A}$  of all operations compatible with all relations in  $\Gamma$ . See Pawel Idziak's talk for an overview of the algebraic approach to CSP.

One of the biggest achievements in the area is a result of Bulatov, which confirms the conjecture in the case that  $\mathbb{A}$  contains all unary relations:

**Theorem 2** (Conservative CSP Dichotomy). *Assume that  $\mathbb{A}$  contains all unary relations over  $A$ . If the corresponding  $\mathbf{A}$  is a Taylor clone, then the problem  $\text{CSP}(\mathbb{A})$  is poly-time, otherwise it is NP-complete.*

A clone  $\mathbf{A}$  is said to be Taylor, if it contains a Taylor operation. This condition has a growing number of equivalent reformulations. For instance, it is known that a finite clone  $\mathbf{A}$  is Taylor, iff the variety generated by  $\mathbf{A}$  omits type 1, iff  $\mathbf{A}$  satisfies some non-trivial idempotent Malcev condition, iff  $\mathbf{A}$  contains a weak near-unanimity operation (or a Siggers operation, or a cyclic operation), ... A variety is called Taylor, if all algebras in the variety have Taylor clones of term operations.

The original proof of Theorem 2 is very technical, complicated and long (the paper has 80 pages). Using a refinement of our work with Marcin Kozik, we are able to give a natural and quite short proof.

In my talk I will present the algebraic core of this new proof – the rectangularity theorem for conservative algebras.

To introduce the rectangularity theorem, we need the concept of an absorbing subalgebra of an algebra.

**Definition 3.** Let  $\mathbf{A}$  be an algebra (or a clone). A subalgebra  $\mathbf{B}$  of  $\mathbf{A}$  is called an *absorbing subalgebra* of  $\mathbf{A}$ , if there exists a term operation  $t$  of  $\mathbf{A}$  such that for every coordinate  $i$  and for every  $b_1, \dots, b_n \in A$ , such that  $b_j \in B$  for all  $j \neq i$ , we have  $t(b_1, \dots, b_n) \in B$ . We denote this fact by  $\mathbf{B} \triangleleft \mathbf{A}$ .

It is not hard to see that the relation  $\triangleleft$  is transitive, therefore the next definition is quite natural.

**Definition 4.** Let  $\mathbf{A}$  be an algebra. A subalgebra  $\mathbf{B}$  of  $\mathbf{A}$  is called a *minimal absorbing subalgebra* (notation  $\mathbf{B} \triangleleft\triangleleft \mathbf{A}$ ), if  $\mathbf{B} \triangleleft \mathbf{A}$  and  $\mathbf{B}$  has no proper absorbing subalgebra.

Recall that an algebra  $\mathbf{A}$  (or a clone) is called conservative, if every operation  $f$  of the algebra (clone) is conservative, that is  $f(a_1, \dots, a_n) \in \{a_1, \dots, a_n\}$  for any  $a_1, \dots, a_n \in A$ . Note that the clone of compatible operations of a relational structure containing all unary relations is conservative.

**Theorem 5** (Rectangularity theorem for conservative algebras). *Let  $\mathbf{A}_1, \dots, \mathbf{A}_n$  be finite conservative algebras from a Taylor variety, let  $\mathbf{B}_i \triangleleft\triangleleft \mathbf{A}_i$ ,  $i = 1, \dots, n$ , let  $R$  be a subdirect subuniverse of  $\mathbf{A}_1 \times \dots \times \mathbf{A}_n$ , let  $S = R \cap (B_1 \times \dots \times B_n)$  and assume  $S \neq \emptyset$ .*

*Consider the following equivalence  $\sim$  on  $\{1, 2, \dots, n\}$ :*

$$i \sim j \text{ iff } \forall (a_1, \dots, a_n) \in R \ (a_i \in B_i \ \& \ a_j \in B_j) \text{ or } (a_i \notin B_i \ \& \ a_j \notin B_j)$$

*Let  $D_1, \dots, D_k$  be a complete list of  $\sim$ -blocks.*

*Then a tuple  $\mathbf{b} = (b_1, \dots, b_n) \in B_1 \times \dots \times B_n$  belongs to  $S$  if and only if the restriction of  $\mathbf{b}$  to  $D_j$  is in the projection of  $R$  to  $D_j$  for every  $j = 1, \dots, k$ .*

*In particular, if all  $\sim$ -blocks are one element, then  $S = B_1 \times \dots \times B_n$ .*