

# **Some properties of strong Galois connections**

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**Michiro Kondo**

**Tokyo Denki University, Japan**

Let  $X$  be a residuated lattice with  $f \dashv g$ . If  $f \dashv g$  satisfies the condition **(sGC1)**:  $g(x) \leq x$ , then

$$\text{(sGC2): } g(f(x) \rightarrow y) = f(x) \rightarrow g(y)$$

$$\iff \left\{ \begin{array}{l} (1) \ g(g(x)) = g(x) \\ (2) \ g(x) \odot g(y) \leq g(x \odot y) \\ (\text{FS}) : f(x \rightarrow y) \leq g(x) \rightarrow f(y) \end{array} \right.$$

## Monadic algebras

monadic Heyting algebras (Bezhanisvili 1998)

monadic MV-algebras (DiNola, Grigolia 2004)

monadic BL-algebras (Grigolia 2006)

monadic  $R^{\ell}$ -monoids (Rachůnek, Švrček 2008)

monadic non-com.  $R^{\ell}$ -monoids (Rachůnek, Šalounova 2013)



monadic residuated lattices from the view  
point of *Galois connection*

$(X, \wedge, \vee, \odot, \rightarrow, 0, 1)$  is a *residuated lattice*, **(RL)** if

- (1)**  $(X, \wedge, \vee, 0, 1)$  is a bounded lattice;
- (2)**  $(X, \odot, 1)$  is a commutative monoid;
- (3)** For all  $x, y, z \in X$ ,

$$x \odot y \leq z \iff x \leq y \rightarrow z$$

**We define  $x' = x \rightarrow 0$ , which is a *negation* in a sense.**

## Proposition 1.

$$(1) \quad 0' = 1, \quad 1' = 0$$

$$(2) \quad x \odot x' = 0$$

$$(3) \quad x \leq y \iff x \rightarrow y = 1$$

$$(4) \quad x \odot (x \rightarrow y) \leq y$$

$$(5) \quad x \leq y \implies x \odot z \leq y \odot z$$

$$(6) \quad x \leq y \implies z \rightarrow x \leq z \rightarrow y, \quad y \rightarrow z \leq x \rightarrow z$$

$$(7) \quad (x \vee y) \odot z = (x \odot z) \vee (y \odot z)$$

$$(8) \quad (x \vee y)' = x' \wedge y'$$

**Let  $(P, \leq)$  be a partially ordered set. A pair of maps  $f, g : P \rightarrow P$  is called a *Galois connection*  $(f \dashv g)$  if**

$$f(x) \leq y \iff x \leq g(y) \quad (x, y \in P)$$

**A Galois connection  $(f, g)$  is *strong*  $(f \dashv_s g)$  if**

**(sGC1)**  $g(x) \leq x$

**(sGC2)**  $g(f(x) \rightarrow y) = f(x) \rightarrow g(y)$

$(X, f, g)$  is a *residuated lattice with strong Galois connection* (**simply, RLsGC**) if  $f \dashv_s g$ .

**Proposition 2.** For RLsGC  $(X, f, g)$ , we have

$$(1) \quad f(0) = 0, g(1) = 1$$

$$(2) \quad x \leq f(x)$$

$$(3) \quad f(g(x)) = g(x), g(f(x)) = f(x)$$

$$(4) \quad f(x) = x \iff g(x) = x$$

$$(5) \quad f(x \odot y) \leq f(x) \odot f(y), \quad g(x) \odot g(y) \leq g(x \odot y)$$

$$(6) \quad f(f(x) \odot f(y)) = f(x) \odot f(y), \\ g(g(x) \odot g(y)) = g(x) \odot g(y)$$

$$(7) \quad g(f(x) \rightarrow y(y)) = f(x) \rightarrow f(y)$$

$$(8) \quad f(f(x) \wedge f(y)) = f(x) \wedge f(y)$$

$$(9) \quad g(f(x) \vee f(y)) = f(x) \vee f(y)$$

**From the above,**

$$X_{fg} = \{x \in X \mid f(x) = x\} = \{x \in X \mid g(x) = x\}$$

**is a subalgebra of  $X$ . A subalgebra  $X_0$  is called *relatively complete* if, for all  $a \in X$ , there exist a minimum and a maximum element of the sets  $\{x \in X_0 \mid a \leq x\}$  and  $\{x \in X_0 \mid x \leq a\}$ , respectively.**

**We denote**

$$f_{X_0}(a) = \min\{x \in X_0 \mid a \leq x\}$$

$$g_{X_0}(a) = \max\{x \in X_0 \mid x \leq a\}$$



**Lemma 3.** For a RLsGC  $(X, f, g)$ ,  $X_{fg}$  is a relatively complete subalgebra.

**Conversely,**

**Proposition 4.** If  $X_0$  is a relatively complete subalgebra, then

$$(1) \text{ (sGC1) } g_{X_0}(x) \leq x$$

$$(2) f_{X_0} \dashv g_{X_0}$$

$$(3) \text{ (sGC2) } g_{X_0}(f_{X_0}(x) \rightarrow y) = f_{X_0}(x) \rightarrow g_{X_0}(y)$$

**From the above,**

**Lemma 5. If  $X_0$  is a relatively complete subalgebra, then  $(X, f_{X_0}, g_{X_0})$  is a RLsGC.**

**Theorem 6. Let  $X$  be a residuated lattice. Then, there exists a strong Galois connection  $f \dashv_s g$   $\iff$  there exists a relatively complete subalgebra of  $X$ .**

$(H, \exists, \forall)$  is called a *monadic Heyting algebra (mHA)* if  $H$  is a Heyting algebra and  $\exists, \forall$  satisfy

$$(H1) \quad \forall x \leq x$$

$$(H2) \quad x \leq \exists x$$

$$(H3) \quad \forall(x \wedge y) = \forall x \wedge \forall y$$

$$(H4) \quad \exists(x \vee y) = \exists x \vee \exists y$$

$$(H5) \quad \forall 1 = 1$$

$$(H6) \quad \exists 0 = 0$$

$$(H7) \quad \forall \exists x = \exists x$$

$$(H8) \quad \exists \forall x = \forall x$$

$$(H9) \quad \exists(\exists x \wedge y) = \exists x \wedge \exists y$$

**Proposition 7.** For every mHA  $(H, \exists, \forall)$ ,  $\exists \dashv_s \forall$ .

**Proposition 8.**  $(H, \exists, \forall)$  is an mHA

$\iff H$  is a RLsGC and  $\wedge = \odot$ .

**Corollary 9.** For a Heyting algebra  $H$ ,

$(H, \exists, \forall)$  is an mHA

$\iff$  there exists a relatively complete subalgebra of  $H$ .

monadic residuated lattice  $(X, \wedge, \vee, \odot, \rightarrow, \exists, \forall, 0, 1)$   
is a residuated lattice with  $\exists, \forall$  satisfying

$$(m1) \quad x \leq \exists x$$

$$(m2) \quad \forall x \leq x$$

$$(m3) \quad \forall(x \rightarrow \exists y) = \exists x \rightarrow \exists y$$

$$(m4) \quad \forall(\exists x \rightarrow y) = \exists x \rightarrow \forall y$$

$$(m5) \quad \forall(x \vee \exists y) = \forall x \vee \exists y$$

$$(m6) \quad \exists \forall x = \forall x$$

$$(m7) \quad \forall \forall x = \forall x$$

$$(m8) \quad \exists(\exists x \odot \exists x) = \exists x \odot \exists x$$

$$(m9) \quad \exists(x \odot x) = \exists x \odot \exists x$$

**Basic result:**  $\exists \dashv_s \forall$ , that is,  $(\exists, \forall)$  is a strong Galois connection.



**Lemma 10.** For a residuated lattice  $X$ ,  
 $(X, \exists, \forall)$  is a monadic residuated lattice

$\iff (X, \exists, \forall)$  is a RLsGC with (m5) and (m9).

**Application: On Kripke frame  $(W, R)$  of modal logic based on CPL,**

$f \dashv g \iff R$  is symmetric

(sGC1) :  $g(x) \leq x \iff R$  is reflexive

(sGC2)  $\iff R$  ?

## Characterization of (sGC2)

$$(\text{sGC2}) : g(f(x) \rightarrow y) = f(x) \rightarrow g(y)$$

**Lemma 11.** Let  $X$  be a residuated lattice with  $f \dashv g$  and **(sGC1)**:  $g(x) \leq x$ . The conditions are equivalent:

**(a) (sGC2):**  $g(f(x) \rightarrow y) = f(x) \rightarrow g(y)$

**(b)**  $g(g(x) \rightarrow y) = g(x) \rightarrow g(y)$

**(c)**  $f(f(x) \odot y) = f(x) \odot f(y)$

**(d)**  $f(g(x) \odot y) = g(x) \odot f(y)$



**Lemma 12.** For any residuated lattice  $X$  with  $f \dashv g$ , the following conditions are equivalent:

**(i)**  $f(x) \rightarrow g(y) \leq g(x \rightarrow y)$

**(ii)**  $f(x) \odot g(y) \leq f(x \odot y)$

**(FS):**  $f(x \rightarrow y) \leq g(x) \rightarrow f(y)$

**Lemma 13.** Let  $X$  be a residuated lattice with  $f \dashv g$  and **(sGC1)**:  $g(x) \leq x$ . Then we have

$$(m3) : g(x \rightarrow f(y)) = f(x) \rightarrow g(y)$$

$$\iff \begin{cases} (1) \ g(g(x)) = g(x) \\ (FS) : f(x \rightarrow y) \leq g(x) \rightarrow f(y) \end{cases}$$

**Lemma 14.** Let  $X$  be a residuated lattice with  $f \dashv g$  and **(sGC1)**:  $g(x) \leq x$ . Then we have

$$\text{(sGC2)} \iff \begin{cases} \text{(m3)} : g(x \rightarrow f(y)) = f(x) \rightarrow g(y) \\ \text{(2)} g(x) \odot g(y) \leq g(x \odot y) \end{cases}$$

Hence we get

**Theorem 15.** Let  $X$  be a residuated lattice with  $f \dashv g$  and **(sGC1)**:  $g(x) \leq x$ .

$$\text{(sGC2)} \iff \left\{ \begin{array}{l} (1) \quad g(g(x)) = g(x) \\ (2) \quad g(x) \odot g(y) \leq g(x \odot y) \\ \text{(FS)} : f(x \rightarrow y) \leq g(x) \rightarrow f(y) \end{array} \right.$$

**Proposition 16.** On any Boolean algebra with Galois connection  $f \dashv g$ ,

$$\text{(FS)} \iff f(x) = (g(x'))'$$

**This means that**

$$\text{CPL} + \{f \dashv g\} + \{gA \rightarrow A\} + \{(\text{sGC2})\}$$

$$= \text{CPL} + \{f \dashv g\} + \{gA \rightarrow A\} \\ + \{gA \rightarrow ggA\} + \{(\text{FS})\}$$

$$= \text{CPL} + \{A \rightarrow gfA\} + \{gA \rightarrow A\} \\ + \{gA \rightarrow ggA\} + \{fA \longleftrightarrow \neg g\neg A\}$$

$$= \text{S5}$$

**Hence, in the case of Heyting algebras (intuitionistic logic),**

$$\text{Int} + \{f \dashv g\} + \{gA \rightarrow A\} + \{(sGC2)\}$$

$$= \text{Int} + \{f \dashv g\} + \{gA \rightarrow A\}$$

$$+ \{gA \rightarrow ggA\} + \{(FS)\}$$

$$= \text{Int} + \{A \rightarrow gfA, fgA \rightarrow A\} + \{gA \rightarrow A\}$$

$$+ \{gA \rightarrow ggA\} + \{(FS)\}$$

$$= \text{IntS5} (= \text{logic of mHA})$$

**Thank you for your attention !!**