## A\&C 6: Complexity theory, NP-completeness.

## Theory.

T6.1 Reductions, polynomial transformations,
T6.2 Complexity classes of decidable problems: P, NP, coNP, NP-hard, NP-complete,
T6.3 Cook-Levin theorem, SAT problem.

## Exercises.

E6.1 Knowing that Cnf-SAT is NP-complete, prove that $k$-SAT is NP-complete for every $k \geq 3$.
E6.2 Show that 3-SAT remains NP-complete, even if each variable appears at most three times.
E6.3 * Show that a variant of 3-SAT, in which every variable appears exactly three times and each clause has exactly three distinct literals, is polynomially solvable.
E6.4 * Show that Clique is NP-complete by a reduction from 3-Sat.
E6.5 Show that Independent Set is NP-complete.
E6.6 Show that Vertex Cover is NP-complete.

## Problems.

Cnf-SAt:
Input: a logic formula in conjunctive normal form, i.e., as a conjunction of clauses, each of which is an alternative of variables or their negations (called literals)
Question: is there an assignment of variables which makes the formula true?
$k$-SAT:
Input: a logic formula in conjunctive normal form, each clause contains at most $k$ literals,
Question: is there an assignment of variables which makes the formula true?
Vertex Cover:
Input: a graph $G$, an integer $k$,
Question: does $G$ contain a set of $k$ vertices, whose removal removes all edges?
Clique:
Input: a graph $G$, an integer $k$,
Question: does $G$ contain a complete subgraph of $k$ vertices?
Independent Set:
Input: a graph $G$, an integer $k$,
Question: does $G$ contain an independent set of $k$ vertices?

