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#### Security Games

- Two asymmetrical players: Defender and Attacker
- Each game is composed of *m* time steps.
- Each player chooses an action to be performed in each time step.
- A player's pure strategy  $\sigma_P$  ( $P \in \{D, A\}$ ) is a sequence of their actions in consecutive time steps:  $\sigma_P = (a_1, a_2, \dots, a_m)$ .
- Defender commits to his/her strategy first.
- Attacker, knowing the Defender's strategy, chooses his/her strategy.
- Defender always commits to a mixed strategy.

Stackelberg equilibrium

**Stackelberg equilibrium**: a pair of players' strategies, for which strategy change by any of players leads to his/her result deterioration.

 $(\pi_D^*, R(\pi_D^*)) \in \Pi_D imes \Pi_A$ 

$$egin{aligned} \pi_D^* &= ext{argmax}_{\pi_D \in \Pi_D} U_D(\pi_D, R(\pi_D)) \ R(\pi_D) &= ext{argmax}_{\pi_A \in \Pi_A} U_A(\pi_D, \pi_A) \end{aligned}$$

 $G \in \{D, A\}$  - players (Defender, Attacker)  $\Pi_G$  - a set of player's G all mixed strategies  $U_G$  - payoff of player G

**Goal:** find optimal Defender's strategy

### Real-life applications



Federal Air Marshal Service



US Coast Guard in Boston Harbor



Los Angeles Airport



Poaching in Uganda



Tickets control in Los Angeles

## Example



#### Example



















### Evolutionary Algorithm for Security Games (EASG)



Żychowski A., Mańdziuk J. *Evolution of Strategies in Sequential Security Games*. Proceedings of the 20th AAMAS conference, pages 1434-1442. 2021

#### EASG - crossover

- Crossover role: combining existing solutions
- Each individual takes part in crossover with crossover rate probability  $p_k$

$$CH_{1-2} = \{ (\sigma_1^1, \frac{p_1^1}{2}), \dots, (\sigma_{l_1}^1, \frac{p_{l_1}^1}{2}), (\sigma_1^2, \frac{p_1^2}{2}), \dots, (\sigma_{l_2}^2, \frac{p_{l_2}^2}{2}) \}$$

• After crossover each pure strategy may be deleted with probability equal to  $(1 - p_i^q)^2$ 



#### EASG - mutation

- Mutation role: introduce some random perturbation to explore new areas of the search space
- Each individual is mutated with mutation rate probability  $p_m$
- Random pure strategy  $\sigma_i^q$  is chosen which is modified starting from the random time step

$$\sigma'_{i}^{q} = (a_{1}, a_{2}, \dots, a_{s-1}, a'_{s}, a'_{s+1}, \dots, a'_{m})$$

EASG mutation - example



#### Mutation enhancements

- **EASG**<sub>n</sub> EASG algorithm with repeated mutation.
- MANPS<sub>1</sub>, MANPS<sub>n</sub> mutation adds new pure strategy a uniformly selected pure strategy is added with a uniformly sampled probability.
- MCP<sub>1</sub>, MCP<sub>n</sub> mutation changes probability a probability of randomly selected pure strategy is uniformly changed.
- MSP<sub>1</sub>, MSP<sub>n</sub> mutation switches probability probabilities of two randomly chosen pure strategies are switched.
- MDPS<sub>1</sub>, MDPS<sub>n</sub> mutation deletes pure strategy -a randomly chosen pure strategy is removed.
- **MCWPS** *mutation changes the weakest pure strategy* mutation is applied only to a pure strategy with the lowest payoff.
- MDWPS mutation deletes the weakest pure strategy pure strategy with the lowest payoff is deleted

### Experimental setup

300 test game instances of 3 types:

- 150 Warehouse Games (WHG)
- 90 Search Games (SEG)
- 60 FlipIt Games (FIG)

30 independent runs for each game instance



Search Games



#### Results

	Defender's payoff			Computation time [s]		
	WHG	SEG	FIG	WHG	SEG	FIG
EASG	0.017	0.108	0.031	152	2534	328
$EASG_n$	0.017	0.135	0.037	1206	21913	3051
$MANPS_1$	0.014	0.059	0.031	156	2548	313
$MANPS_n$	0.016	<u>0.139</u>	0.036	1366	21892	2988
$MCP_1$	0.015	0.074	0.030	148	2422	336
$MCP_n$	0.016	<u>0.131</u>	0.037	1285	22651	3008
$MSP_1$	0.013	0.099	0.024	156	2583	316
$MSP_n$	0.016	0.108	0.037	1332	21447	2931
$MDPS_1$	0.013	0.052	0.029	147	2620	313
$MDPS_n$	0.013	0.053	0.026	1283	22026	2900
MCWPS	0.013	0.046	0.030	148	2612	321
MDWPS	0.008	0.058	0.018	139	2361	299

The average Defender's payoff and the computation time for various mutation operators. The best results are **bolded**. Results that are better than the baseline version of the algorithm (EASG) are <u>underlined</u>. In cases where the difference between the baseline version (EASG) and a given variation is statistically significant the result is highlighted with a gray background.

#### Conclusions

- Repetition of mutation operation leads to improvement of SSGs outcomes, though at the expense of significant increase in computation time.
- The proposed modifications offer a viable alternative to the base EASG formulation for cases when computational cost is less important.

#### Thank you





# **MANPS** - *mutation adds new pure strategy* - a uniformly selected pure strategy is added with a uniformly sampled probability



# **MCP** - *mutation changes probability* - a probability of randomly selected pure strategy is uniformly changed



**MSP** - *mutation switches probability* - probabilities of two randomly chosen pure strategies are switched



## **MDPS** - *mutation deletes pure strategy* -a randomly chosen pure strategy is removed



**MCWPS -** *mutation changes the weakest pure strategy* - mutation is applied only to a pure strategy with the lowest payoff



**MDWPS** - *mutation deletes the weakest pure strategy* - pure strategy with the lowest payoff is deleted













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